# On representing the positive semidefinite cone using the second-order cone 

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#### Abstract

The second-order cone plays an important role in convex optimization and has strong expressive abilities despite its apparent simplicity. Second-order cone formulations can also be solved more efficiently than semidefinite programming problems in general. We consider the following question, posed by Lewis and Glineur, Parrilo, Saunderson: is it possible to express the general positive semidefinite cone using second-order cones? We provide a negative answer to this question and show that the $3 \times 3$ positive semidefinite cone does not admit any second-order cone representation. In fact we show that the slice consisting of $3 \times 3$ positive semidefinite Hankel matrices does not admit a second-order cone representation. Our proof relies on exhibiting a sequence of submatrices of the slack matrix of the $3 \times 3$ positive semidefinite cone whose "second-order cone rank" grows to infinity.


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## 1 Introduction

Let $\mathcal{Q} \subset \mathbb{R}^{3}$ denote the three-dimensional second-order cone (also known as the "ice-cream" cone or the Lorentz cone):

$$
\mathcal{Q}=\left\{(x, t) \in \mathbb{R}^{2} \times \mathbb{R}:\|x\| \leq t\right\}
$$

[^0]It is known that $\mathcal{Q}$ is linearly isomorphic to the cone of $2 \times 2$ real symmetric positive semidefinite matrices. Indeed we have:

$$
\left(x_{1}, x_{2}, t\right) \in \mathcal{Q} \Longleftrightarrow\left[\begin{array}{cc}
t-x_{1} & x_{2}  \tag{1}\\
x_{2} & t+x_{1}
\end{array}\right] \succeq 0
$$

Despite its apparent simplicity the second-order cone $\mathcal{Q}$ has strong expressive abilities and allows us to represent various convex constraints that go beyond "simple quadratic constraints". For example it can be used to express geometric means $\left(x \mapsto \prod_{i=1}^{n} x_{i}^{p_{i}}\right.$ where $p_{i} \geq 0$, rational, and $\sum_{i=1}^{n} p_{i}=1$ ), $\ell_{p}$-norm constraints, multifocal ellipses (see e.g., [11, Equation (3.5)]), robust counterparts of linear programs, etc. We refer the reader to [4, Section 3.3] for more details.

Given this strong expressive ability one may wonder whether the general positive semidefinite cone can be represented using $\mathcal{Q}$. This question was posed in particular by Adrian Lewis (personal communication) and Glineur, Parrilo and Saunderson [7]. In this paper we show that this is not possible, even for the $3 \times 3$ positive semidefinite cone. To make things precise we use the language of lifts (or extended formulations), see [8]. We denote by $\mathcal{Q}^{k}$ the Cartesian product of $k$ copies of $\mathcal{Q}$ :

$$
\mathcal{Q}^{k}=\mathcal{Q} \times \cdots \times \mathcal{Q} \quad(k \text { copies }) .
$$

A linear slice of $\mathcal{Q}^{k}$ is an intersection of $\mathcal{Q}^{k}$ with a linear subspace. We say that a convex cone $K \subset \mathbb{R}^{m}$ has a second-order cone lift of size $k$ (or simply $\mathcal{Q}^{k}$-lift) if it can be written as the projection of a slice of $\mathcal{Q}^{k}$, i.e.:

$$
\begin{equation*}
K=\pi\left(\mathcal{Q}^{k} \cap L\right) \tag{2}
\end{equation*}
$$

where $\pi: \mathbb{R}^{3 k} \rightarrow \mathbb{R}^{m}$ is a linear map and $L$ is a linear subspace of $\mathbb{R}^{3 k}$. Let $\mathbf{S}_{+}^{n}$ be the cone of $n \times n$ real symmetric positive semidefinite matrices. In this paper we prove:

Theorem 1 The cone $\mathbf{S}_{+}^{3}$ does not admit any $\mathcal{Q}^{k}$-lift for any finite $k$.
Actually our proof allows us to show that the slice of $\mathbf{S}_{+}^{3}$ consisting of Hankel matrices does not admit any second-order representation (see Sect. 4 for details). Note that higher-dimensional second order cones of the form

$$
\left\{(x, t) \in \mathbb{R}^{n} \times t:\|x\| \leq t\right\}
$$

where $n \geq 3$ can be represented using the three-dimensional cone $\mathcal{Q}$, see e.g., [5, Section 2]. Thus Theorem 1 also rules out any representation of $\mathbf{S}_{+}^{3}$ using the higherdimensional second-order cones. Moreover since $\mathbf{S}_{+}^{3}$ appears as a slice of higher-order positive semidefinite cones Theorem 1 also shows that one cannot represent $\mathbf{S}_{+}^{n}$, for $n \geq 3$ using second-order cones.

## 2 Preliminaries

The paper [8] introduced a general methodology to prove existence or nonexistence of lifts in terms of the slack matrix of a cone. In this section we review some of the definitions and results from this paper, and introduce the notion of a second-order cone factorization and the second-order cone rank.

Let $E$ be a Euclidean space with inner product $\langle\cdot, \cdot\rangle$ and let $K \subseteq E$ be a cone. The dual cone $K^{*}$ is defined as:

$$
K^{*}=\{x \in E:\langle x, y\rangle \geq 0 \quad \forall y \in K\} .
$$

We also denote by $\operatorname{ext}(K)$ the extreme rays of a cone $K$. The notion of slack matrix plays a fundamental role in the study of lifts.

Definition 1 (Slack matrix) The slack matrix of a cone $K$, denoted $S_{K}$, is a (potentially infinite) matrix where columns are indexed by extreme rays of $K$, and rows are indexed by extreme rays of $K^{*}$ (the dual of $K$ ) and where the $(x, y)$ entry is given by:

$$
\begin{equation*}
S_{K}[x, y]=\langle x, y\rangle \quad \forall(x, y) \in \operatorname{ext}\left(K^{*}\right) \times \operatorname{ext}(K) \tag{3}
\end{equation*}
$$

Note that, by definition of dual cone, all the entries of $S_{K}$ are nonnegative. Also note that an element $x \in \operatorname{ext}\left(K^{*}\right)$ (and similarly $y \in \operatorname{ext}(K)$ ) is only defined up to a positive multiple. Any choice of scaling gives a valid slack matrix of $K$ and the properties of $S_{K}$ that we are interested in will be independent of the scaling chosen.

The existence/nonexistence of a second-order cone lift for a convex cone $K$ will depend on whether $S_{K}$ admits a certain second-order cone factorization which we now define.

Definition 2 ( $\mathcal{Q}^{k}$-factorization and second-order cone rank) Let $S \in \mathbb{R}^{|I| \times|J|}$ be a matrix with nonnegative entries. We say that $S$ has a $\mathcal{Q}^{k}$-factorization if there exist vectors $a_{i} \in \mathcal{Q}^{k}$ for $i \in I$ and $b_{j} \in \mathcal{Q}^{k}$ for $j \in J$ such that $S[i, j]=\left\langle a_{i}, b_{j}\right\rangle$ for all $i \in I$ and $j \in J$. The smallest $k$ for which such a factorization exists will be denoted $\operatorname{rank}_{\mathrm{soc}}(S)$.

Remark 1 Recall that for any $a, b \in \mathcal{Q}$ we have $\langle a, b\rangle \geq 0$. This means that any matrix with a second-order cone factorization is elementwise nonnegative.

Remark 2 It is important to note that the second-order cone rank of any matrix $S$ can be equivalently expressed as the smallest $k$ such that $S$ admits a decomposition

$$
\begin{equation*}
S=M_{1}+\cdots+M_{k} \tag{4}
\end{equation*}
$$

where $\operatorname{rank}_{\mathrm{soc}}\left(M_{l}\right)=1$ for each $l=1, \ldots, k$ (i.e., each $M_{l}$ has a factorization $M_{l}[i, j]=\left\langle a_{i}, b_{j}\right\rangle$ where $\left.a_{i}, b_{j} \in \mathcal{Q}\right)$. This simply follows from the fact that $\mathcal{Q}^{k}$ is the Cartesian product of $k$ copies of $\mathcal{Q}$.

We now state the result from [8] that we will need.

Theorem 2 (Existence of a lift, special case of [8]) Let $K$ be a convex cone. If $K$ has a $\mathcal{Q}^{k}$-lift then its slack matrix $S_{K}$ has a $\mathcal{Q}^{k}$-factorization.

This theorem can actually be turned into an if and only if condition under mild conditions on $K$ (e.g., $K$ is proper), see [8], but we have only stated here the direction that we will need.

The cone $\mathbf{S}_{+}^{3}$ In this paper we are interested in the cone $K=\mathbf{S}_{+}^{3}$ of real symmetric $3 \times 3$ positive semidefinite matrices. The extreme rays of $\mathbf{S}_{+}^{3}$ are rank-one matrices of the form $x x^{T}$ where $x \in \mathbb{R}^{3}$. Also $\mathbf{S}_{+}^{3}$ is self-dual, i.e., $\left(\mathbf{S}_{+}^{3}\right)^{*}=\mathbf{S}_{+}^{3}$. The slack matrix of $\mathbf{S}_{+}^{3}$ thus has its rows and columns indexed by three-dimensional vectors and

$$
\begin{equation*}
S_{\mathbf{S}_{+}^{3}}[x, y]=\left\langle x x^{T}, y y^{T}\right\rangle=\left(x^{T} y\right)^{2} \quad \forall(x, y) \in \mathbb{R}^{3} \times \mathbb{R}^{3} . \tag{5}
\end{equation*}
$$

In order to prove that $\mathbf{S}_{+}^{3}$ does not admit a second-order representation, we will show that its slack matrix does not admit any $\mathcal{Q}^{k}$-factorization for any finite $k$. In fact we will exhibit a sequence $\left(A_{n}\right)$ of submatrices of $S_{\mathbf{S}_{+}^{3}}$ where $\operatorname{rank}_{\text {soc }}\left(A_{n}\right)$ grows to $+\infty$ as $n \rightarrow+\infty$.

Before introducing this sequence of matrices we record the following simple (known) proposition concerning orthogonal vectors in the cone $\mathcal{Q}$ which will be useful later.

Proposition 1 Let a, $b_{1}, b_{2} \in \mathcal{Q}$ nonzero and assume that $\left\langle a, b_{1}\right\rangle=\left\langle a, b_{2}\right\rangle=0$. Then $b_{1}$ and $b_{2}$ are collinear.

Proof This is easy to see geometrically by visualizing the "ice cream" cone. We include a proof for completeness: let $a=\left(a^{\prime}, t\right) \in \mathbb{R}^{2} \times \mathbb{R}$ and $b_{i}=\left(b_{i}^{\prime}, s_{i}\right) \in \mathbb{R}^{2} \times \mathbb{R}$ where $\left\|a^{\prime}\right\| \leq t$ and $\left\|b_{i}^{\prime}\right\| \leq s_{i}$. Note that for $i=1,2$ we have $0=\left\langle a, b_{i}\right\rangle=\left\langle a^{\prime}, b_{i}^{\prime}\right\rangle+t s_{i} \geq$ $-\left\|a^{\prime}\right\|\left\|b_{i}^{\prime}\right\|+t s_{i} \geq 0$ where in the first inequality we used Cauchy-Schwarz and in the second inequality we used the definition of the second-order cone. It thus follows that all the inequalities must be equalities: by the equality case in Cauchy-Schwarz we must have that $b_{i}^{\prime}=\alpha_{i} a^{\prime}$ for some constant $\alpha_{i}<0$ and we must also have $t=\left\|a^{\prime}\right\|$ and $s_{i}=\left\|b_{i}^{\prime}\right\|$. Thus we get that $b_{i}=\left(\alpha_{i} a^{\prime},\left|\alpha_{i}\right|\left\|a^{\prime}\right\|\right)=\left|\alpha_{i}\right|\left(-a^{\prime},\left\|a^{\prime}\right\|\right)$. This shows that $b_{1}$ and $b_{2}$ are both collinear to the same vector $\left(-a^{\prime},\left\|a^{\prime}\right\|\right)$ and thus completes the proof.

## 3 Proof of Theorem 1

A sequence of matrices We now define our sequence $A_{n}$ of submatrices of the slack matrix of $\mathbf{S}_{+}^{3}$. For any integer $i$ define the vector

$$
\begin{equation*}
v_{i}=\left(1, i, i^{2}\right) \in \mathbb{R}^{3} \tag{6}
\end{equation*}
$$

Note that this sequence of vectors satisfies the following:

$$
\begin{equation*}
\text { For all distinct integers } i_{1}, i_{2}, i_{3} \operatorname{det}\left(v_{i_{1}}, v_{i_{2}}, v_{i_{3}}\right) \neq 0 \tag{7}
\end{equation*}
$$

Our matrix $A_{n}$ has size $\binom{n}{2} \times n$ and is defined as follows (rows are indexed by 2 -subsets of $[n]$ and columns are indexed by $[n]$ ):

$$
\begin{align*}
A_{n}\left[\left\{i_{1}, i_{2}\right\}, j\right]: & =\left(\left(v_{i_{1}} \times v_{i_{2}}\right)^{T} v_{j}\right)^{2} \\
& =\operatorname{det}\left(v_{i_{1}}, v_{i_{2}}, v_{j}\right)^{2} \quad \forall\left\{i_{1}, i_{2}\right\} \in\binom{[n]}{2}, \forall j \in[n] \tag{8}
\end{align*}
$$

where $\times$ denotes the cross-product of three-dimensional vectors. It is clear from the definition of $A_{n}$ that it is a submatrix of the slack matrix of $\mathbf{S}_{+}^{3}$. Note that the sparsity pattern of $A_{n}$ satisfies the following:

$$
\begin{array}{ll}
A_{n}[e, j]=0 & \text { if } j \in e  \tag{9}\\
A_{n}[e, j]>0 & \text { otherwise }
\end{array} \quad e \in\binom{[n]}{2}, j \in[n] .
$$

Also note that $A_{n}$ satisfies the following important recursive property: for any subset $C$ of $[n]$ of size $n_{0}$ the submatrix $A_{n}\left[\binom{C}{2}, C\right]$ has the same sparsity pattern as $A_{n_{0}}$ (up to relabeling of rows and columns). In our main theorem we will show that the second-order cone rank of $A_{n}$ grows to infinity with $n$.

Remark 3 (Geometric interpretation of (9)) The property (9) of the matrices $A_{n}$ will be the key to prove a lower bound on their second-order cone rank. Geometrically, the property (9) reflects a certain 2-neighborliness property of the extreme rays ${ }^{1} \operatorname{ext}\left(\mathbf{S}_{+}^{3}\right)$ of $\mathbf{S}_{+}^{3}$ : for any two distinct extreme rays $x x^{T}$ and $y y^{T}$ of $\mathbf{S}_{+}^{3}$, there is a supporting hyperplane $H$ to $\mathbf{S}_{+}^{3}$ that touches $\operatorname{ext}\left(\mathbf{S}_{+}^{3}\right)$ precisely at $x x^{T}$ and $y y^{T}$. This 2-neighborliness property turns out to be the key geometric obstruction for the existence of second-order cone lifts for $\mathbf{S}_{+}^{3}$.

Covering numbers Our analysis of the matrix $A_{n}$ will only rely on its sparsity pattern. Given two matrices $A$ and $B$ of the same size we write $A \stackrel{\text { supp }}{=} B$ if $A$ and $B$ have the same support (i.e., $A_{i j}=0$ if and only if $B_{i j}=0$ for all $i, j$ ). We now define a combinatorial analogue of the second-order cone rank:

Definition 3 Given a nonnegative matrix $A$, we define the soc-covering number of $A$, denoted $\operatorname{cov}_{\mathrm{soc}}(A)$ to be the smallest number $k$ of matrices $M_{1}, \ldots, M_{k}$ with $\operatorname{rank}_{\mathrm{soc}}\left(M_{l}\right)=1$ for $l=1, \ldots, k$ that are needed to cover the nonzero entries of $A$, i.e., such that

$$
\begin{equation*}
A \stackrel{\text { supp }}{=} M_{1}+\cdots+M_{k} . \tag{10}
\end{equation*}
$$

Proposition 2 For any nonnegative matrix $A$ we have $\operatorname{rank}_{\mathrm{soc}}(A) \geq \operatorname{cov}_{\mathrm{soc}}(A)$.
Proof This follows immediately from Remark 2 concerning $\operatorname{rank}_{\text {soc }}$ and the definition of $\operatorname{cov}_{\text {soc }}$.

[^1]A simple but crucial fact concerning soc-coverings that we will use is the following: in any soc-covering of $A$ of the form (10), each matrix $M_{l}$ must satisfy $M_{l}[i, j]=0$ whenever $A[i, j]=0$. This is because the matrices $M_{1}, \ldots, M_{k}$ are all entrywise nonnegative.

We are now ready to state our main result.
Theorem 3 Consider a sequence $\left(A_{n}\right)$ of matrices of sparsity pattern given in (9). Then for any $n_{0} \geq 2$ we have $\operatorname{cov}_{\mathrm{soc}}\left(A_{3 n_{0}^{2}}\right) \geq \operatorname{cov}_{\mathrm{soc}}\left(A_{n_{0}}\right)+1$. As a consequence $\operatorname{cov}_{\text {soc }}\left(A_{n}\right) \rightarrow+\infty$ when $n \rightarrow+\infty$.

The proof of our theorem rests on a key lemma concerning the sparsity pattern of any term in a soc-covering of $A_{n}$.

Lemma 1 (Main) Let $n$ be such that $n \geq 3 n_{0}^{2}$ for some $n_{0} \geq 2$. Assume $M \in \mathbb{R}\binom{n}{2} \times n$ satisfies $\operatorname{rank}_{\mathrm{soc}}(M)=1$ and $M[e, j]=0$ for all $e \in\binom{n}{2}$ and $j \in[n]$ such that $j \in e$. Then there is a subset $C$ of $[n]$ of size at least $n_{0}$ such that the submatrix $M\left[\binom{C}{2}, C\right]$ is identically zero.

Before proving this lemma, we show how this lemma can be used to easily prove Theorem 3.

Proof of Theorem 3 Let $n=3 n_{0}^{2}$ and consider a soc-covering of $A_{n} \stackrel{\text { supp }}{=} M_{1}+\cdots+$ $M_{r}$ of size $r=\operatorname{cov}_{\text {soc }}\left(A_{n}\right)$ (note that we have of course $r \geq 1$ since $A_{n}$ is not identically zero). By Lemma 1 there is a subset $C$ of $[n]$ of size $n_{0}$ such that $M_{1}\left[\binom{C}{2}, C\right]=0$. It thus follows that we have $A_{n}\left[\binom{C}{2}, C\right] \stackrel{\text { supp }}{=} M_{2}\left[\binom{C}{2}, C\right]+\cdots+M_{r}\left[\binom{C}{2}, C\right]$. Also note that $A_{n}\left[\binom{C}{2}, C\right] \stackrel{\text { supp }}{=} A_{n_{0}}$. It thus follows that $A_{n_{0}}$ has a soc-covering of size $r-1$ and thus $\operatorname{cov}_{\mathrm{soc}}\left(A_{n_{0}}\right) \leq \operatorname{cov}_{\mathrm{soc}}\left(A_{3 n_{0}^{2}}\right)-1$. This completes the proof.
For completeness we show how Theorem 1 follows directly from Theorem 3.
Proof of Theorem 1 Since for any $n \geq 1, A_{n}$ is a submatrix of the slack matrix of $\mathbf{S}_{+}^{3}$, Theorem 3 shows that the slack matrix of $\mathbf{S}_{+}^{3}$ does not admit any $\mathcal{Q}^{k}$-factorization for finite $k$. This shows, via Theorem 2 , that $\mathbf{S}_{+}^{3}$ does not have a $\mathcal{Q}^{k}$-lift for any finite $k$.

The rest of the section is devoted to the proof of Lemma 1.
Proof of Lemma 1 Let $M \in \mathbb{R}^{\binom{n}{2} \times n}$ and assume that $M$ has a factorization $M_{e, j}=$ $\left\langle a_{e}, b_{j}\right\rangle$ where $a_{e}, b_{j} \in \mathcal{Q}$ for all $e \in\binom{[n]}{2}$ and $j \in[n]$, and that $M_{e, j}=0$ whenever $j \in e$.

Let $E_{0}:=\left\{e \in\binom{[n]}{2}: a_{e}=0\right\}$ be the set of rows of $M$ that are identically zero and let $E_{1}=\binom{[n]}{2} \backslash E_{0}$. Similarly for the columns we let $S_{0}:=\left\{j \in[n]: b_{j}=0\right\}$ and $S_{1}=[n] \backslash S_{0}$.

In the next lemma we use the sparsity pattern of $A_{n}$ together with Proposition 1 to infer additional properties on the sparsity pattern of $M$.

Lemma 2 Let $C$ be a connected component of the graph with vertex set $S_{1}$ and edge set $E_{1}\left(S_{1}\right)$ (where $E_{1}\left(S_{1}\right)$ consists of elements in $E_{1}$ that connect only elements of $S_{1}$ ). Then necessarily $M\left[\binom{C}{2}, C\right]=0$.

Proof We first show using Proposition 1 that all the vectors $\left\{b_{j}\right\}_{j \in C}$ are necessarily collinear. Let $j_{1}, j_{2} \in S_{1}$ such that $e=\left\{j_{1}, j_{2}\right\} \in E_{1}$. Note that since $M_{e, j_{1}}=$ $M_{e, j_{2}}=0$ then we have, by Proposition 1 that $b_{j_{1}}$ and $b_{j_{2}}$ are collinear. It is easy to see thus now that if $j_{1}$ and $j_{2}$ are connected by a path in the graph $\left(S_{1}, E_{1}\left(S_{1}\right)\right)$ then $b_{j_{1}}$ and $b_{j_{2}}$ must be collinear.

We thus get that all the columns of $M$ indexed by $C$ must be proportional to each other, and so they must have the same sparsity pattern. Now let $e \in\binom{C}{2}$. If $a_{e}=0$ then $M[e, C]=0$ since the entire row indexed by $e$ is zero. Otherwise if $a_{e} \neq 0$ let $e=\left\{j_{1}, j_{2}\right\}$ with $j_{1}, j_{2} \in C$. Since, by assumption, $M_{e, j_{1}}=0$ it follows that for any $j \in C$ we must have $M_{e, j}=0$, i.e., $M[e, C]=0$. This is true for any $e \in\binom{C}{2}$ thus we get that $M\left[\binom{C}{2}, C\right]=0$.

To complete the proof of Lemma 1 assume that $n \geq 3 n_{0}^{2}$ for some $n_{0} \geq 2$. We need to show that there is a subset $C$ of $[n]$ of size at least $n_{0}$ such that $\left.M\left[\begin{array}{c}C \\ 2\end{array}\right), C\right]=0$.

First note that if the graph $\left(S_{1}, E_{1}\left(S_{1}\right)\right)$ has a connected component of size at least $n_{0}$ then we are done by Lemma 2. Also note that if $S_{0}$ has size at least $n_{0}$ we are also done because all the columns indexed by $S_{0}$ are identically zero by definition.

In the rest of the proof we will thus assume that $\left|S_{0}\right|<n_{0}$ and that the connected components of $\left(S_{1}, E_{1}\left(S_{1}\right)\right)$ all have size $<n_{0}$. We will show in this case that $E_{0}$ necessarily contains a clique of size at least $n_{0}$ (i.e., a subset of the form $\binom{C}{2}$ where $|C| \geq n_{0}$ ) and this will prove our claim since all the rows in $E_{0}$ are identically zero by definition. The intuition is as follows: the assumption that $\left|S_{0}\right|<n_{0}$ and that the connected components of ( $S_{1}, E_{1}\left(S_{1}\right)$ ) have size $<n_{0}$ mean that the graph $\left(S_{1}, E_{1}\left(S_{1}\right)\right)$ is very sparse. In particular this means that $E_{1}$ has to be small which means that $E_{0}=E_{1}^{c}$ must be large and thus it must contain a large clique.

More precisely, to show that $E_{1}$ is small note that it consists of those edges that are either in $E_{1}\left(S_{1}\right)$ or, otherwise, they must have at least one node in $S_{1}^{c}=S_{0}$. Thus we get that

$$
\left|E_{1}\right| \leq\left|E_{1}\left(S_{1}\right)\right|+\left|S_{0}\right|(n-1) \leq\left|E_{1}\left(S_{1}\right)\right|+\left(n_{0}-1\right)(n-1) .
$$

where in the second inequality we used the fact that $\left|S_{0}\right|<n_{0}$. Also since the connected components of $\left(S_{1}, E_{1}\left(S_{1}\right)\right)$ all have size $<n_{0}$ it is not difficult to show that $\left|E_{1}\left(S_{1}\right)\right|<$ $n_{0} n / 2$ (indeed if we let $x_{1}, \ldots, x_{k}$ be the size of each connected component we have $\left.\left|E_{1}\left(S_{1}\right)\right| \leq \frac{1}{2} \sum_{i=1}^{k} x_{i}^{2}<\frac{1}{2} \sum_{i=1}^{k} n_{0} x_{i} \leq \frac{1}{2} n_{0} n\right)$. Thus we get that

$$
\left|E_{1}\right| \leq \frac{n_{0} n}{2}+\left(n_{0}-1\right)(n-1) \leq\left(\frac{3}{2} n_{0}-1\right) n
$$

Thus this means, since $E_{0}=\binom{[n]}{2} \backslash E_{1}$ :

$$
\left|E_{0}\right| \geq\binom{ n}{2}-\left(\frac{3}{2} n_{0}-1\right) n>\frac{n^{2}}{2}-\frac{3}{2} n_{0} n
$$

We now invoke a result of Turán to show that $E_{0}$ must contain a clique of size at least $n_{0}$ :

Theorem 4 (Turán, see e.g., [2]) Any graph on $n$ vertices with more than $\left(1-\frac{1}{k}\right) \frac{n^{2}}{2}$ edges contains a clique of size $k+1$.

By taking $k=n_{0}-1$ we see that $E_{0}$ contains a clique of size $n_{0}$ if

$$
\frac{n^{2}}{2}-\frac{3}{2} n_{0} n>\left(1-\frac{1}{n_{0}-1}\right) \frac{n^{2}}{2}
$$

This simplifies into

$$
n>3 n_{0}\left(n_{0}-1\right)
$$

which is true for $n \geq 3 n_{0}^{2}$.

## 4 Slices of the $3 \times 3$ positive semidefinite cone

Hankel slice The proof given in the previous section actually shows the following more general statement.

Theorem 5 Let $\mathcal{H}$ denote the cone of $3 \times 3$ positive semidefinite Hankel matrices:

$$
\mathcal{H}=\left\{X \in \mathbf{S}_{+}^{3}: X_{13}=X_{22}\right\} .
$$

Assume $K$ is a convex cone that is "sandwiched" between $\mathcal{H}$ and $\mathbf{S}_{+}^{3}$, i.e., $\mathcal{H} \subseteq K \subseteq$ $\mathbf{S}_{+}^{3}$. Then $K$ does not have a second-order cone representation.

Proof The proof follows from the observation that the matrices $A_{n}$ considered in Sect. 3 (see Eq. (8)) are actually submatrices of the generalized slack matrix of the pair of nested cones $\left(\mathcal{H}, \mathbf{S}_{+}^{3}\right)$, the definition of which we now recall (see e.g., [9, Definition 6]): The generalized slack matrix of a pair of convex cones ( $K_{1}, K_{2}$ ) with $K_{1} \subseteq K_{2}$ is a matrix whose rows are indexed by ext $\left(K_{2}^{*}\right)$ (the valid linear inequalities of $K_{2}$ ) and its columns indexed by ext $\left(K_{1}\right)$ and is defined by

$$
S_{K_{1}, K_{2}}[x, y]=\langle x, y\rangle \quad x \in \operatorname{ext}\left(K_{2}^{*}\right), y \in \operatorname{ext}\left(K_{1}\right)
$$

When $K_{1}=K_{2}$ this is precisely the slack matrix of $K_{1}=K_{2}$. The following theorem is a generalization of Theorem 2 and can be proved using very similar arguments (see e.g., [9, Proposition 7]).

Theorem 6 (Generalization of Theorem 2 to nested cones) Let $K_{1}, K_{2}$ be two convex cones with $K_{1} \subseteq K_{2}$, and assume there exists a convex cone $K$ with a $\mathcal{Q}^{k}$-lift such that $K_{1} \subseteq K \subseteq K_{2}$. Then $S_{K_{1}, K_{2}}$ has a $\mathcal{Q}^{k}$-factorization.

The main observation is to see that the vector $v_{i}$ defined in (6) satisfies $v_{i} v_{i}^{T} \in \mathcal{H}$, and so this shows that $A_{n}$ defined in Equation (8) is a submatrix of the generalized slack matrix of the pair $\left(\mathcal{H}, \mathbf{S}_{+}^{3}\right)$. Since $\operatorname{rank}_{\text {soc }}\left(A_{n}\right)$ grows to infinity with $n$, this gives the desired result.

The dual cone of $\mathcal{H}$ is the cone of nonnegative quartic polynomials on the real line (see e.g., [3, Section 3.5]). It thus follows that the latter is also not second-order cone representable using the second-order cone. More generally we can prove:

Corollary 1 Let $\Sigma_{n, 2 d}$ be the cone of polynomials in $n$ variables of degree at most $2 d$ that are sums of squares. Then $\Sigma_{n, 2 d}$ is not second-order cone representable except in the case $(n, 2 d)=(1,2)$.

For the proof we recall that in the cases $n=1$ (univariate polynomials) and $2 d=2$ (quadratic polynomials), nonnegative polynomials are sums of squares.

Proof For $2 d=2, \Sigma_{n, 2 d}$ is the cone of nonnegative quadratic polynomials in $n$ variables. By homogenization, this cone is linearly isomorphic to $\mathbf{S}_{+}^{n+1}$, the cone of nonnegative quadratic forms in $n+1$ variables. By Theorem 1, this shows that $\Sigma_{n, 2}$ is not second-order cone representable for $n \geq 2$. The case $(n, 2 d)=(1,2)$ is clearly second-order cone representable because $\mathbf{S}_{+}^{2}$ is linearly isomorphic to the second-order cone.

If $2 d \geq 4$ then the cone of nonnegative quartic polynomials on the real line can be obtained as a section of $\Sigma_{n, 2 d}$ by setting to zero the coefficients of some appropriate monomials. This shows that $\Sigma_{n, 2 d}$ is not second-order cone representable when $2 d \geq$ 4.

Other slices of $S_{+}^{3}$ that are second-order cone representable There are certain slices of $\mathbf{S}_{+}^{3}$ of codimension 1 that are, on the other hand, known to admit a second-order cone representation. For example the following second-order cone representation of the slice $\left\{X \in \mathbf{S}_{+}^{3}: X_{11}=X_{22}\right\}$ appears in [7]:

$$
\begin{align*}
{\left[\begin{array}{ccc}
t & a & b \\
a & t & c \\
b & c & s
\end{array}\right] \succeq 0 \Longleftrightarrow } & \exists u, v \in \mathbb{R} \text { s.t. }\left[\begin{array}{cc}
t+a & b+c \\
b+c & u
\end{array}\right] \succeq 0, \\
& {\left[\begin{array}{cc}
t-a & b-c \\
b-c & v
\end{array}\right] \succeq 0, \quad u+v=2 s . } \tag{11}
\end{align*}
$$

(The $2 \times 2$ positive semidefinite constraints can be converted to second-order cone constraints using (1)). To see why (11) holds note that by applying a congruence transformation by $\frac{1}{\sqrt{2}}\left[\begin{array}{ccc}1 & 1 & 0 \\ 1 & -1 & 0 \\ 0 & 0 & 2\end{array}\right]$ on the $3 \times 3$ matrix on the left-hand side of (11) we get that

$$
\left[\begin{array}{ccc}
t & a & b \\
a & t & c \\
b & c & s
\end{array}\right] \succeq 0 \Longleftrightarrow\left[\begin{array}{ccc}
t+a & 0 & b+c \\
0 & t-a & b-c \\
b+c & b-c & 2 s
\end{array}\right] \succeq 0
$$

The latter matrix has an arrow structure and thus using results on the decomposition of matrices with chordal sparsity pattern $[1,6,10]$ we get the decomposition (11).

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[^1]:    ${ }^{1}$ In fact here we only work with the extreme rays $\left\{v_{n} v_{n}^{T}: n \in \mathbb{N}\right\}$, see Sect. 4 for the implication of this.

